

Data l'espressione:

$$\frac{|\eta_{+-}|^2}{|\eta_{00}|^2} = \frac{\left[BR(K_L \rightarrow \pi^+ \pi^-) \right]}{\left[BR(K_S \rightarrow \pi^+ \pi^-) \right]} \cong 1 + 6 \Re e\left(\frac{\epsilon'}{\epsilon}\right)$$

Dimostrare che :

$$\delta \Re e\left(\frac{\epsilon'}{\epsilon}\right)_{stat} = \frac{1}{6} \frac{1}{\sqrt{(2/3)N_L^0}}$$

con N_L^0 numero di conteggi $K_L \rightarrow \pi^0 \pi^0$.

In quale approssimazione vale la formula?

La relazione fra $BR_{S,L}^{\pm,0}$ e $N_{S,L}^{\pm,0}$ è data da:

$$N_{S,L}^{\pm,0} = N_{S,L}^{\pm,0}(obs) - Bck_{S,L}^{\pm,0} = N_{KK} \cdot \rho_{S,L}(tag) \cdot BR_{S,L}^{\pm,0} \cdot \langle \rho_{S,L}^{\pm,0} \rangle \cdot \iint_{FV} g(l-l') I(l) dl dl'$$

dove:

- $N_{S,L}^{\pm,0}(obs)$ è il numero effettivamente osservato di decadimenti $\pi^+ \pi^- , \pi^0 \pi^0$;
- N_{KK} è il numero totale di coppie K_S, K_L prodotte;
- $\rho_{S,L}(tag)$ è l'efficienza di identificazione;
- $BR_{S,L}^{\pm,0}$ è il "branching ratio" corrispondente al decadimento $K_{S,L} \rightarrow \pi^+ \pi^- , \pi^0 \pi^0$;
- $\langle \rho_{S,L}^{\pm,0} \rangle$ è l'efficienza media di rivelazione dei decadimenti $K_{S,L} \rightarrow \pi^+ \pi^- , \pi^0 \pi^0$;
- $\iint_{FV} g(l-l') I(l) dl dl'$ rappresenta la convoluzione dell'intensità $I(l) = e^{-\eta l_{SL}}$ dei decadimenti con la risoluzione sperimentale $g(l-l')$ sul cammino di decadimento l, integrata sul volume fiduciale del rivelatore.

- $Bck_{S,L}^{\pm,0}$ è il contributo degli eventi di fondo.

$$N_L^0 = \underbrace{\sigma_{e^+ e^- \rightarrow \phi}}_{\sigma_{e^+ e^- \rightarrow \phi}} \cdot \underbrace{\mathcal{L}}_{\int L dt} \cdot \underbrace{0.66}_{\rho_L(tag)} \cdot \underbrace{0.34}_{BR(\phi \rightarrow K_S K_L)} \cdot \underbrace{10^{-3}}_{BR_L^0} \cdot \underbrace{(e^{-30350} - e^{-150350})}_{fiducial volume}$$

Data una luminosità integrata di $\mathcal{L}=10^4 \text{ pb}^{-1}$ qual'è il fattore di reiezione del fondo $K_L \rightarrow 3\pi^0$ (sul segnale $K_L \rightarrow 2\pi^0$) necessario per avere un errore su $\text{Re}(\epsilon'/\epsilon) < 3 \times 10^{-4}$ assumendo di conoscere il fondo con una precisione del 20%?

Soluzione

1) A parita' di luminosita', i decadimenti del K_L in due pioni sono molto piu' rari di quelli corrispondenti del K_S .

Quindi l'errore statistico e' dominato dai conteggi sul K_L .

Trascurando l'incertezza sulle efficienze e la sottrazione del fondo, considerando che $N^0_L \sim 2 N^\pm_L$ si ottiene la formula data.

2) Si generalizza la formula precedente ad includere l'errore derivante dalla sottrazione del fondo nei conteggi N^0_L (attenzione: tenere separati i contributi di N^0_L e N^\pm_L). Considerando che

$$N(K_L \rightarrow 3\pi^0) \sim 200 N^0_L \text{ e che } N_B = N(K_L \rightarrow 3\pi^0)/R$$

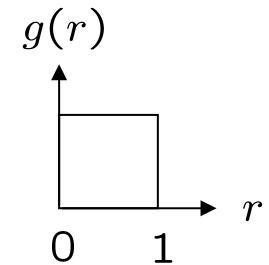
date le condizioni indicate si ottiene per il fattore di reiezione: $R > 25800$.

The Monte Carlo method

What it is: a numerical technique for calculating probabilities and related quantities using sequences of random numbers.

The usual steps:

- (1) Generate sequence r_1, r_2, \dots, r_m uniform in $[0, 1]$.
- (2) Use this to produce another sequence x_1, x_2, \dots, x_n distributed according to some pdf $f(x)$ in which we're interested (x can be a vector).
- (3) Use the x values to estimate some property of $f(x)$, e.g., fraction of x values with $a < x < b$ gives $\int_a^b f(x) dx$.
→ MC calculation = integration (at least formally)



MC generated values = ‘simulated data’

→ use for testing statistical procedures

Random number generators

Goal: generate uniformly distributed values in $[0, 1]$.

Toss coin for e.g. 32 bit number... (too tiring).

→ ‘random number generator’

= computer algorithm to generate r_1, r_2, \dots, r_n .

Example: multiplicative linear congruential generator (MLCG)

$$n_{i+1} = (a n_i) \bmod m , \quad \text{where}$$

n_i = integer

a = multiplier

m = modulus

n_0 = seed (initial value)

N.B. mod = modulus (remainder), e.g. $27 \bmod 5 = 2$.

This rule produces a sequence of numbers n_0, n_1, \dots

Random number generators (2)

The sequence is (unfortunately) periodic!

Example (see Brandt Ch 4): $a = 3, m = 7, n_0 = 1$

$$n_1 = (3 \cdot 1) \bmod 7 = 3$$

$$n_2 = (3 \cdot 3) \bmod 7 = 2$$

$$n_3 = (3 \cdot 2) \bmod 7 = 6$$

$$n_4 = (3 \cdot 6) \bmod 7 = 4$$

$$n_5 = (3 \cdot 4) \bmod 7 = 5$$

$$n_6 = (3 \cdot 5) \bmod 7 = 1 \quad \leftarrow \text{sequence repeats}$$

Choose a, m to obtain long period (maximum = $m - 1$); m usually close to the largest integer that can be represented in the computer.

Only use a subset of a single period of the sequence.

Random number generators (3)

$r_i = n_i/m$ are in $[0, 1]$ but are they ‘random’?

Choose a, m so that the r_i pass various tests of randomness:

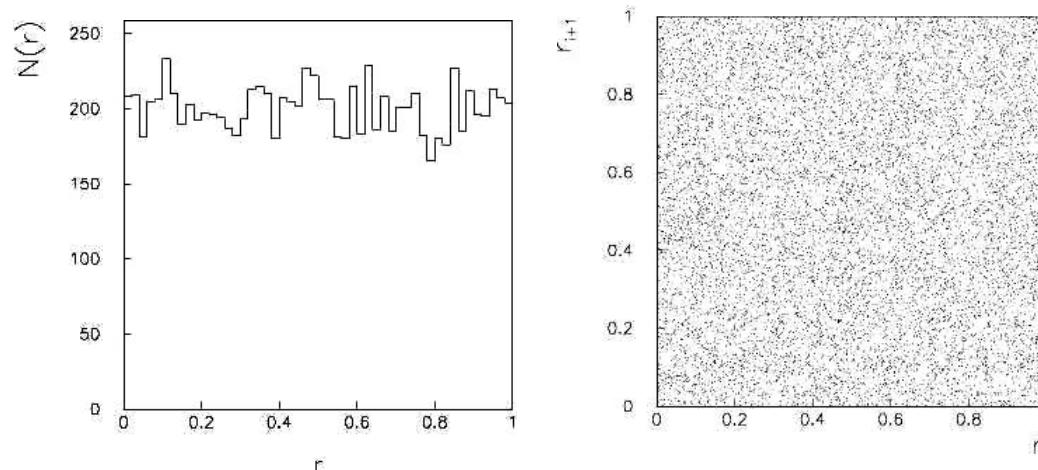
uniform distribution in $[0, 1]$,

all values independent (no correlations between pairs),

e.g. L'Ecuyer, Commun. ACM 31 (1988) 742 suggests

$$a = 40692$$

$$m = 2147483399$$

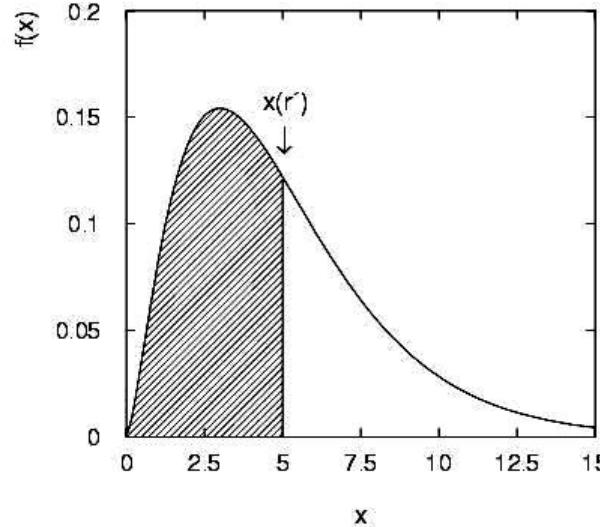
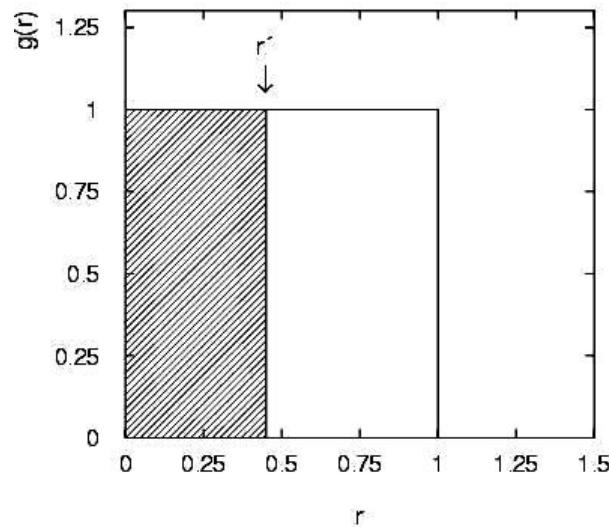


Far better generators available, e.g. TRandom3, based on Mersenne twister algorithm, period = $2^{19937} - 1$ (a “Mersenne prime”).

See F. James, Comp. Phys. Comm. 60 (1990) 111; Brandt Ch. 4

The transformation method

Given r_1, r_2, \dots, r_n uniform in $[0, 1]$, find x_1, x_2, \dots, x_n that follow $f(x)$ by finding a suitable transformation $x(r)$.



Require: $P(r \leq r') = P(x \leq x(r'))$

i.e. $\int_{-\infty}^{r'} g(r) dr = r' = \int_{-\infty}^{x(r')} f(x') dx' = F(x(r'))$

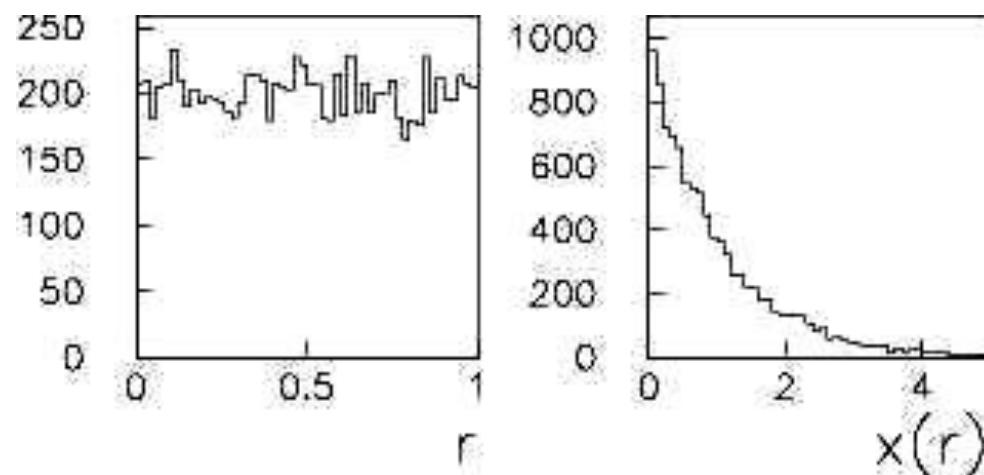
That is, set $F(x) = r$ and solve for $x(r)$.

Example of the transformation method

Exponential pdf: $f(x; \xi) = \frac{1}{\xi} e^{-x/\xi}$ ($x \geq 0$)

Set $\int_0^x \frac{1}{\xi} e^{-x'/\xi} dx' = r$ and solve for $x(r)$.

→ $x(r) = -\xi \ln(1 - r)$ ($x(r) = -\xi \ln r$ works too.)



1. Generation of random,
Pseudo-random numbers
2. Random variable r uniformly distributed between 0 and 1
3. Sampling of a discrete random variable

Example:

A discrete random variable x with 3 values, x_1, x_2, x_3 with probabilities P_1, P_2 and P_3 respectively ($\sum P_i = 1$).

Extract $y=r$

if $0 < y < P_1 \Rightarrow x = x_1$

if $P_1 < y < (P_1 + P_2) \Rightarrow x = x_2$

if $(P_1 + P_2) < y < 1 \Rightarrow x = x_3$

4. Sampling of a continuous random variable x with arbitrary pdf $f(x)$

Extract $y=r$

$x=F^{-1}(y)$ with $y=F(x)=\int_0^x f(x')dx'$

Example:

$f(x)=1/(b-a) \Rightarrow x=a+(b-a)r$

$f(\theta)=\sin\theta/2 \Rightarrow \cos\theta=1-2r \Rightarrow \theta=\arccos(1-2r)$

$f(x)=\mu \exp(-\mu x) \Rightarrow x=-\ln(1-r)/\mu \Rightarrow x=-\ln(r)/\mu$